

**Welcome to INF0216:
Knowledge Graphs
Spring 2022**

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Session 10: Reasoning about KGs (DL)

- Themes:
 - description logic
 - decision problems
 - OWL DL
 - Manchester OWL-syntax



Readings

- Material at <http://wiki.uib.no/info216> (cursory):
 - <http://www.w3.org/TR/owl2-primer/>
 - show: Turtle and Manchester syntax
 - hide: other syntaxes
 - Description Logic Handbook:
 - Chapter 1: Nardi & Brachman: Introduction to Description Logics
 - Chapter 2: Baader & Nutt: Formal Description Logics (gets hard)

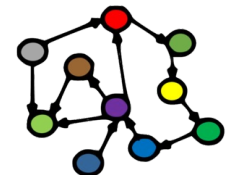


THE KNOWLEDGE GRAPH
COOKBOOK
RECIPES THAT WORK

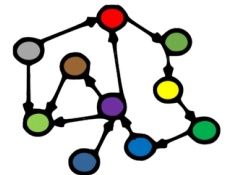


ANDREAS BLUMAUER
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1st edition, 2020



Description Logic (DL)

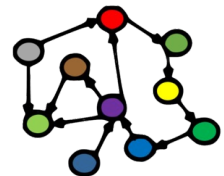


Relationship to other logics

- *Proposition logics* are about *statements (propositions)*:
 "Martha is a Woman" \Leftarrow
 "Martha is Human" \wedge "Martha is Female"
- (First order) *predicate logics* are about *predicates* and *objects*:
 - $\forall \mathbf{x} . (\text{Woman}(\mathbf{x}) \Leftrightarrow \text{Human}(\mathbf{x}) \wedge \text{Female}(\mathbf{x}))$
- *Description logics* are about *concepts*:
 - $\text{Woman} \doteq \text{Human} \sqcap \text{Female}$
 - ...and also about *roles* and *individuals*
- There are many other logic systems:
 - *modal logics*: necessarily \square , possibly \diamond
 - *temporal logics*: always \square , sometimes \diamond , next time \circ

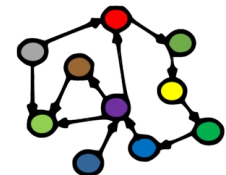
Description logics

- Description Logic (DL)
 - a simple *fragment* of predicate logic
 - ...or, rather, a *family of such fragments*
 - not very *expressive* (“uttrykkskraftig”)
 - but (can have) *good decision problems*, i.e.,
 - it answers many *decision problems* (rather) quickly
- Suitable for describing *concepts* (“begreper”)
 - formal basis for *OWL DL*
 - can be used to:
 - describe *concepts* and their *roles* (“**Tbox**”)
 - describe *roles* and their relations (“**Rbox**”)
 - describe *individuals* and their *roles* (“**ABox**”)



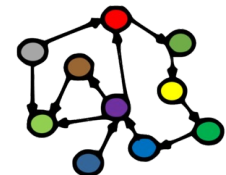
Definition of concepts (“begreper”)

- **Woman** \doteq **Human** \sqcap **Female**
- **Man** \doteq **Human** \sqcap \neg **Woman**
- **Parent** \doteq **Mother** \sqcup **Father**
 - **concepts**: **Human, Female, Woman...**
 - **definition**: \doteq
 - **conjunction** (and): \sqcap
 - **disjunction** (or): \sqcup
 - **negation** (not): \neg
 - **nested expressions**: ()
- **Childless** \doteq ..using **Human** and **Parent**..



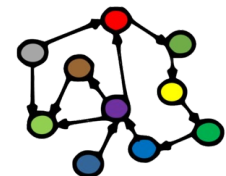
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 - **disjunction** (or): \sqcup
 - **negation** (not): \neg
 - **nested expressions**: ()
- **Childless** \doteq **Human** \sqcap \neg **Parent**



Types of concepts (“begreper”)

- **Woman** \doteq **Human** \sqcap **Female**
- **Man** \doteq **Human** \sqcap \neg **Woman**
- **Parent** \doteq **Mother** \sqcup **Father**
 - atomic (or basic, primitive) concepts:
Human, Female, Woman...
 - only used on the r.h.s. of definitions
 - concept expressions (complex concepts):
 \neg **Woman, Human** \sqcap **Female...**
 - only used on the r.h.s. of definitions
 - defined (and named) concepts:
Woman, Man...
 - defined on the l.h.s. of definitions



Atomic and defined concepts

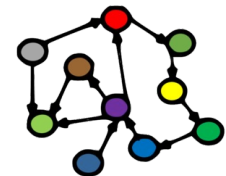
- **Atomic (or basic) concepts**
 - given, always named
 - cannot appear on the l.h.s. of a \doteq definition
 - correspond to simple OWL-NamedClasses
- **Concept expressions**
 - defined in terms of other concepts (and roles)
 - correspond to complex OWL-Classes
- Defined concepts can also be **named**
 - must appear on the l.h.s. of a \doteq definition
 - **concept_name** \doteq **concept_expression**
- ...similar distinction between atomic and defined **roles**



Roles

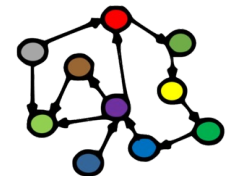
An atomic
(or basic) role

- **Mother** \doteq **Female** \sqcap \exists hasChild. \top
- **Bachelor** \doteq **Male** \sqcap $\neg\exists$ hasSpouse. \top
- **Uncle** \doteq **Male** \sqcap \exists hasSibling.Parent
 - roles: hasChild, hasSibling...
 - universal concept ("top"): \top
 - existential restriction: \exists
- **Grandparent** \doteq ..using Human, hasChild, Parent..
- **Grandparent** \doteq ..using only Human, hasChild..
- **Uncle** \doteq ..using Male, hasSibling, hasChild..



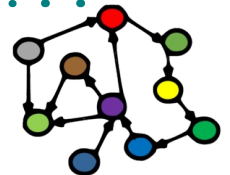
Roles

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- **Bachelor** \doteq **Male** \sqcap $\neg\exists$ **hasSpouse**. \top
- **Uncle** \doteq **Male** \sqcap \exists **hasSibling**.**Parent**
 - `roles: hasChild, hasSibling...`
 - `universal concept ("top"): \top`
 - `existential restriction: \exists`
- **Grandparent** \doteq **Human** \sqcap \exists **hasChild**.**Parent**
- **Grandparent** \doteq ..using only `Human, hasChild..`
- **Uncle** \doteq ..using `Male, hasSibling, hasChild..`



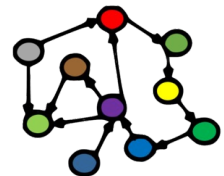
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 - `roles`: hasChild, hasSibling...
 - universal concept ("top"): \top
 - existential restriction: \exists
- **Grandparent** \doteq **Human** \sqcap \exists hasChild.Parent
- **Grandparent** \doteq **Human** \sqcap \exists hasChild. \exists hasChild. \top
- **Uncle** \doteq ...using Male, hasSibling, hasChild....



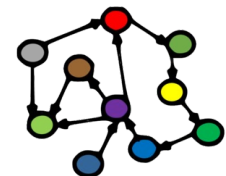
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 - `roles: hasChild, hasSibling...`
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- **Grandparent** \doteq **Human** \sqcap \exists hasChild. \exists hasChild. \top
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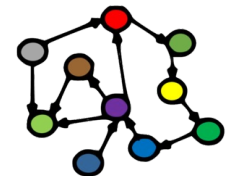
Null concept

- **Male** \sqcap **Female** $\sqsubseteq \perp$
 - null concept (“bottom”): \perp
 - subsumption (sub concept): \sqsubseteq
- \sqsubseteq is used for *subsumption axioms*
 - or: containment / specialisation axioms
- \doteq is used for *definitions* (or just \equiv)
 - \equiv is also used for *equivalence axioms*
- Note the use of $\dots \sqsubseteq \perp$ (“subsumption of bottom”) to say that something is not the case



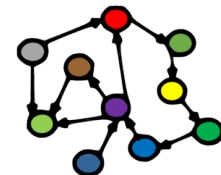
Null concept

- **Male** \sqsupset **Female** $\sqsubseteq \perp$
 - null concept (“bottom”): \perp
 - subsumption (sub concept): \sqsubseteq
- \sqsubseteq is used for *subsumption axioms*
 - or: containment / specialisation axioms
- \doteq is used for *definitions* (or just \equiv)
 - \equiv is also used for *equivalence axioms*
- Note the use of $\dots \sqsubseteq \perp$ (“subsumption of bottom”) to say that something is not the case
- *This was our first proper axiom!*
 - so far we have just *defined* concepts
 - we have not used them in proper *axioms*



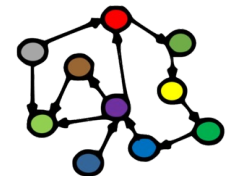
Axioms

- \doteq is used for *definitions*
- \equiv is used for *equivalence axioms*
 - and sometimes for *definitions* too...
- Axioms are equivalences or subsumptions:
 - *subsumption axioms* (\sqsubseteq):
 - composite concept (role) expressions on both sides
 - *equivalence axioms* (\equiv):
 - composite concept (role) expressions on both sides
 - corresponds to: $C \sqsubseteq D, D \sqsubseteq C$
- *expression* $\sqsubseteq \perp$ (“subsumption of bottom”) is used to say that something is *not* the case



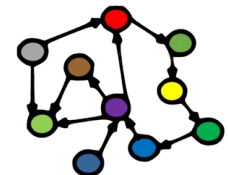
More role definitions

- **HappyFather** \doteq **Father** \sqcap \forall **hasChild.HappyPerson**
 - universal restriction: \forall
- **MotherOfOne** \doteq **Mother** \sqcap **=1 hasChild.** \top
- **Polygamist** \doteq **≥ 3 hasSpouse.** \top
 - number restrictions: $=, \geq, \leq$
- **Narsissist** \doteq \exists **hasLoveFor.****Self**
 - self references: **Self**
- **MassMurderer** \doteq ...using **hasKilled, Human...**



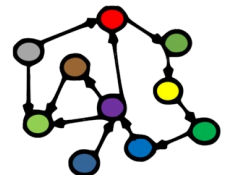
More uses of roles

- **HappyFather** \doteq **Father** \sqcap \forall **hasChild.HappyPerson**
 - universal restriction: \forall
- **MotherOfOne** \doteq **Mother** \sqcap **=1 hasChild.** \top
- **Polygamist** \doteq ≥ 3 **hasSpouse.** \top
 - number restrictions: $=, \geq, \leq$
- **Narsissist** \doteq \exists **hasLoveFor.****Self**
 - self references: **Self**
- **MassMurderer** \doteq ≥ 4 **hasKilled.Human**



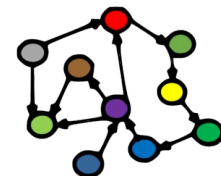
Inverse and transitive roles

- $\text{Child} \doteq \text{Human} \sqcap \exists \text{hasChild}^- . \top$
- $\text{hasParent} \doteq \text{hasChild}^-$
- $\text{hasSibling} \doteq \text{hasSibling}^-$
- $\text{BlueBlood} \doteq \forall \text{hasParent}^* . \text{BlueBlood}$
 - inverse role: hasChild^-
 - symmetric role: hasSibling^-
 - transitive role: hasParent^*
- $\text{Niece} \doteq ..\text{Woman}, \text{hasChild}, \text{hasSibling}..$



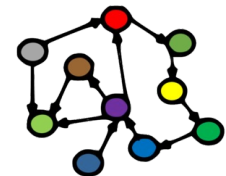
Inverse and transitive roles

- $\text{Child} \doteq \text{Human} \sqcap \exists \text{hasChild}^- . \top$
- $\text{hasParent} \doteq \text{hasChild}^-$
- $\text{hasSibling} \doteq \text{hasSibling}^-$
- $\text{BlueBlood} \doteq \forall \text{hasParent}^* . \text{BlueBlood}$
 - inverse role: hasChild^-
 - symmetric role: hasSibling^-
 - transitive role: hasParent^*
- $\text{Niece} \doteq \text{Woman} \sqcap \exists \text{hasChild}^- . \text{hasSibling} . \top$
- *We just started to define roles!*
 - until now, we have only defined *concepts*



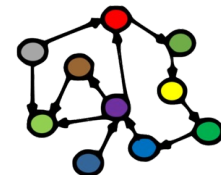
Composite roles

- Similar to composite concepts, e.g.:
 - **hasUncle** \doteq **hasParent** \circ **hasBrother**
 - **hasLovedChild** \doteq **hasChild** \sqcap **hasLoveFor**
 - **hasBrother** \doteq (**hasSibling** | **Male**)
- Not always supported by reasoning engines
 - they can have “bad decision problems”
 - i.e., they compute slowly or intractably
 - ...with some exceptions
- **hasDaughter** \doteq ..using **hasChild**, **Female**..



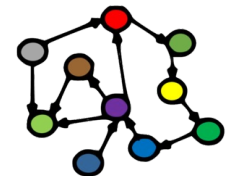
Composite roles

- Similar to composite concepts, e.g.:
 - **hasUncle** \doteq **hasParent** \circ **hasBrother**
 - **hasLovedChild** \doteq **hasChild** \sqcap **hasLoveFor**
 - **hasBrother** \doteq (**hasSibling** | **Male**)
- Not always supported by reasoning engines
 - they can have “bad decision problems”
 - i.e., they compute slowly or intractably
 - ...with some exceptions
- **hasDaughter** \doteq (**hasChild** | **Female**)



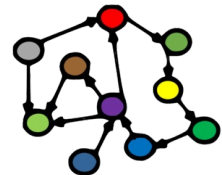
TBox

- *Terminology box* (TBox):
 - a collection of definitions
 - *definition axioms* (\doteq):
 - $\text{concept_name} \doteq \text{concept_expression}$
 - defined and named concept on the l.h.s.
 - complex concept expression on the r.h.s
 - *defined names*
 - must appear on the l.h.s. of some \doteq definition
 - *atomic (basic, primitive) names*
 - can only appear on the r.h.s. of \doteq definitions



Acyclic, definitional TBox

- **Source** \doteq \exists **hasSource⁻.Content**
- **TrustedContent** \doteq \exists **hasSource.TrustedSource**
- **VerifiedContent** \doteq \exists **verifiedBy.FactChecker**
- **DebunkedContent** \doteq \exists **debunkedBy.FactChecker**
- **UnreliableSource** \doteq \exists **hasSource⁻.DebunkedContent**
- **VerifyingSource** \doteq \exists **hasSource⁻.VerifiedContent**
 \sqcap \forall **hasSource⁻.VerifiedContent**

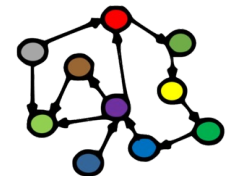


Acyclic, definitional TBox

- **Source** \doteq
 - **TrustedContent** \doteq
 - **VerifiedContent** \doteq
 - **DebunkedContent** \doteq
 - **UnreliableSource** \doteq
 - **VerifyingSource** \doteq
- Defined concepts \square

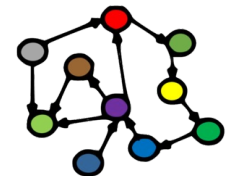
- \exists **hasSource⁻.Content** Concept expressions of atomic concepts
- \exists **hasSource.TrustedSource**
- \exists **verifiedBy.FactChecker**
- \exists **debunkedBy.FactChecker**
- \exists **hasSource⁻.DebunkedContent**
- \exists **hasSource⁻.VerifiedContent**
- \forall **hasSource⁻.VerifiedContent**

Acyclic and unequivocal!



TBox

- **Acyclicity**: no cyclic definitions in the TBox
- **Unequivocality**: each named defined term is only used on the l.h.s. of a single definition
- **Concept expansion**:
 - every concept can be written as an expression of only atomic concepts
 - algorithm:
 - start with the expression that defines the concept
 - recursively replace all the defined concepts used in the expression with their definitions
 - halt when only atomic concepts remain



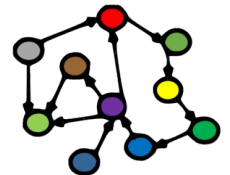
Expanded definitional TBox

- **Source** \doteq \exists **hasSource⁻.Content**
- **TrustedContent** \doteq \exists **hasSource.TrustedSource**
- **VerifiedContent** \doteq \exists **verifiedBy.FactChecker**
- **DebunkedContent** \doteq \exists **debunkedBy.FactChecker**
- **UnreliableSource** \doteq \exists **hasSource⁻.**
 \exists **debunkedBy.FactChecker**
- **VerifyingSource** \doteq \exists **hasSource⁻.**
 \exists **verifiedBy.FactChecker**
 $\sqcap \forall$ **hasSource⁻.**
 \exists **verifiedBy.FactChecker**

*Only basic concepts on
the right hand sides!*

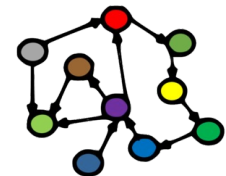
RBox

- *Role box* (RBox):
 - a collection of definitions *of roles*
 - otherwise similar to TBoxes:
 - atomic (basic, primitive) roles
 - role expressions
 - named defined roles
 - role expansion
 - not always necessary (i.e., only atomic roles)



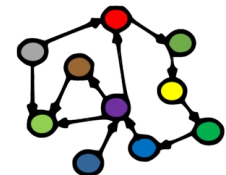
ABox

- So far definitions of concepts and roles (*TBox*, *RBox*)
- Also two types of axioms about individuals (*ABox*):
 - *class assertion* (using a *concept*):
Märtha : Female \sqcap Royal
 - *role assertion* (using a *role*):
<Märtha, EmmaTallulah> : hasChild
<Märtha, HaakonMagnus> : hasBrother
- A TBox + an ABBox (+ possibly an RBox) constitute a *knowledge base* (\mathcal{K}):
 - concepts, roles in the *TBox* (aka “the tags”)
 - roles in the *RBox* (also “tags”)
 - individuals, roles in the *ABBox* (“the tagged data”)



Syntaxes differ a bit...

- So far definitions of concepts and roles (*TBox*, *RBox*)
- Also two types of axioms about individuals (*ABox*):
 - *class assertion* (using a *concept*):
`Female(Märtha) , (Female \sqcap Royal) (Märtha)`
 - *role assertion* (using a *role*):
`hasChild(Märtha, EmmaTallulah)`
`hasBrother(Märtha, HaakonMagnus)`
- A TBox + an ABox (+ possibly an RBox) constitute a *knowledge base* (\mathcal{K}):
 - concepts, roles in the *TBox* (aka “the tags”)
 - roles in the *RBox* (also “tags”)
 - individuals, roles in the *ABox* (“the tagged data”)



Summary of axioms

- Terminology axioms (TBox):

– subsumptions: $C \sqsubseteq D$

– equivalences: $C \equiv D$

corresponds to: $C \sqsubseteq D, D \sqsubseteq C$

C and D are *expressions*!

- Role axioms (RBox)

- Individual assertion axioms (in the ABox):

– class assertions: $a : C$

– role assertions: $\langle a, b \rangle : R$

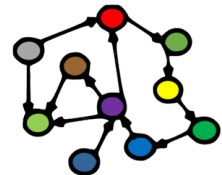
a and b are *individuals*.
R is a *role*!

- Knowledge base $\mathcal{K} = (\mathcal{T}, \mathcal{A})$ or $\mathcal{K} = (\mathcal{T}, \mathcal{R}, \mathcal{A})$

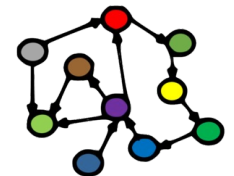
– TBox: \mathcal{T}

RBox: \mathcal{R}

ABox: \mathcal{A}

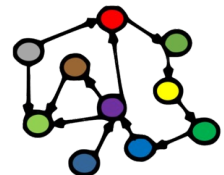


Decision Problems



Reasoning over knowledge bases

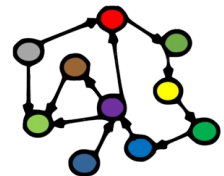
- *What more can we do with ontologies?*
- For example:
 - given a *source ontology* that describes media content along with its sources and trustworthiness
 - we can *answer questions* like, e.g.:
 - is trusted content a type of content?
 - can content be both verified and debunked?
 - is all verified content trusted?
 - *competency questions* are what we want an ontology to answer
 - *DL offers a clear and compact way of representing and reasoning about questions such as these!*



Decision problems

- A computational problem with a yes/no answer, e.g.
 - is C *subsumed* by D: $\mathcal{K} \models C \sqsubseteq D$?
 - are C and D *consistent*: $\mathcal{K} \models a : (C \sqcap D)$?
 - does a *belong* to C: $\mathcal{K} \models a : C$?
 - is a *R-related* to b: $\mathcal{K} \models \langle a, b \rangle : R$?
- Given a knowledge base \mathcal{K} , reasoning engines are designed to give yes / no answer
 - ...but not all decision problems are *decidable*
 - ...or have tractable *complexity*
 - *depends on the expressions used!*

C and D are classes,
a and b are individuals.
R is a role!



Decision problems for concepts

- Four important decision problems for concepts:

- consistency:

can an individual **a** exist so that

$$\mathcal{T} \models \mathbf{a} : \mathbf{C}$$

- subsumption:

$$\mathcal{T} \models \mathbf{C} \sqsubseteq \mathbf{D}$$

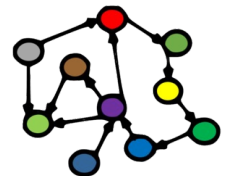
- equivalence:

$$\mathcal{T} \models \mathbf{C} \equiv \mathbf{D}, \text{ also written } \mathbf{C} \equiv_{\mathcal{T}} \mathbf{D},$$

- disjunction:

$$\mathcal{T} \models \mathbf{C} \sqcap \mathbf{D} \sqsubseteq \perp$$

- \mathcal{T} can always be *emptied*, by expanding all its concepts



Decision problems for concepts

- *All four can be reduced to subsumption or consistency!*

– consistency:

$$\mathcal{T} \models \mathbf{a:C} \leftrightarrow \mathcal{T} \not\models \mathbf{C} \sqsubseteq \perp$$

$$\mathcal{T} \Box \mathbf{a:C} \leftrightarrow \mathcal{T} \models \mathbf{C} \sqsubseteq \perp$$

– subsumption:

$$\mathcal{T} \models \mathbf{C} \sqsubseteq \mathbf{D} \leftrightarrow \mathcal{T} \models \mathbf{C} \sqcap \neg \mathbf{D} \sqsubseteq \perp$$

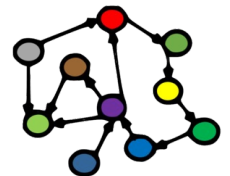
– equivalence:

$$\mathcal{T} \models \mathbf{C} \equiv \mathbf{D} \leftrightarrow \mathcal{T} \models \mathbf{C} \sqsubseteq \mathbf{D}, \mathbf{D} \sqsubseteq \mathbf{C}$$

– disjunction:

$$\mathcal{T} \models \mathbf{C} \sqcup \mathbf{D} \sqsubseteq \perp$$

- \mathcal{T} can always be *emptied*, by expanding all its concepts



Decision problems for individuals

- Decision problems for individuals and roles:
 - instance checking:
 - is individual **a** member of class/concept **C**?
 - $\mathcal{A} \models \mathbf{a}:\mathbf{C}$ $\neq \mathcal{A} \sqcap \neg(\mathbf{a}:\mathbf{C})$
 - role checking:
 - is individual **a** **R**-related to individual **b**?
 - $\mathcal{A} \models \langle \mathbf{a}, \mathbf{b} \rangle : \mathbf{R}$ $\neq \mathcal{A} \sqcap \neg(\langle \mathbf{a}, \mathbf{b} \rangle : \mathbf{R})$
 - classifications (not yes/no):
 - to which classes/concepts does **a** belong?
 - all individuals of class/concept **C**?
- *Everything boils down to consistency checking for ABoxes*
 - ...under certain (rather weak) conditions



Decision problems in practice

- Description logic is implemented by *reasoning engines/OWL reasoner*
 - e.g., *HerMiT and Pellet*
 - distinct from *inference engines*, such as OWL-RL
- Protegé-OWL
 - comes with HerMiT, more plugins can be installed
- Owlready2 (an OWL programming API built around)
 - comes with HerMiT and Pellet, HerMiT is default
- Solves decision problems, e.g.,
 - classify individuals
 - find subclass relationships (subsumptions)
 - find unsatisfiable classes (concepts)
 - detect inconsistencies

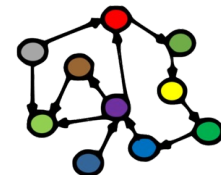
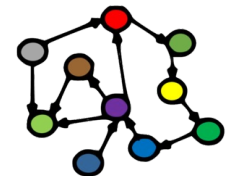


Tableau algorithm

- A simple reasoning procedure
- Tests satisfiability of a concept C_0
 - C_0 is possible expanded
 - negation normal form (NNF)
- Starts with ABox $A_0 = \{ C_0(x) \}$
- Applies transformation rules that preserve consistency
- Halt when not more rules can be applied
 - ...and halt a branch that contains a contradiction
- If all possible branches contain contradictions:
 - C_0 is unsatisfiable
- C_0 is satisfiable otherwise



The \rightarrow_{\sqcap} -rule

Condition: \mathcal{A} contains $(C_1 \sqcap C_2)(x)$, but it does not contain both $C_1(x)$ and $C_2(x)$.

Action: $\mathcal{A}' = \mathcal{A} \cup \{C_1(x), C_2(x)\}$.

The \rightarrow_{\sqcup} -rule

Condition: \mathcal{A} contains $(C_1 \sqcup C_2)(x)$, but neither $C_1(x)$ nor $C_2(x)$.

Action: $\mathcal{A}' = \mathcal{A} \cup \{C_1(x)\}$, $\mathcal{A}'' = \mathcal{A} \cup \{C_2(x)\}$.

The \rightarrow_{\exists} -rule

Condition: \mathcal{A} contains $(\exists R.C)(x)$, but there is no individual name z such that $C(z)$ and $R(x, z)$ are in \mathcal{A} .

Action: $\mathcal{A}' = \mathcal{A} \cup \{C(y), R(x, y)\}$ where y is an individual name not occurring in \mathcal{A} .

The \rightarrow_{\forall} -rule

Condition: \mathcal{A} contains $(\forall R.C)(x)$ and $R(x, y)$, but it does not contain $C(y)$.

Action: $\mathcal{A}' = \mathcal{A} \cup \{C(y)\}$.

The \rightarrow_{\geq} -rule

Condition: \mathcal{A} contains $(\geq n R)(x)$, and there are no individual names z_1, \dots, z_n such that $R(x, z_i)$ ($1 \leq i \leq n$) and $z_i \neq z_j$ ($1 \leq i < j \leq n$) are contained in \mathcal{A} .

Action: $\mathcal{A}' = \mathcal{A} \cup \{R(x, y_i) \mid 1 \leq i \leq n\} \cup \{y_i \neq y_j \mid 1 \leq i < j \leq n\}$, where y_1, \dots, y_n are distinct individual names not occurring in \mathcal{A} .

The \rightarrow_{\leq} -rule

Condition: \mathcal{A} contains distinct individual names y_1, \dots, y_{n+1} such that $(\leq n R)(x)$ and $R(x, y_1), \dots, R(x, y_{n+1})$ are in \mathcal{A} , and $y_i \neq y_j$ is not in \mathcal{A} for some $i \neq j$.

Action: For each pair y_i, y_j such that $i > j$ and $y_i \neq y_j$ is not in \mathcal{A} , the ABox $\mathcal{A}_{i,j} = [y_i/y_j]\mathcal{A}$ is obtained from \mathcal{A} by replacing each occurrence of y_i by y_j .

Next week:
Formal ontologies
(OWL-DL)